



UNIVERSITÀ DEGLI STUDI DELL' AQUILA

Distributed Systems: Mid-term Evaluation

Tuesday, November 24th, 2015 – Prof. Guido Proietti

Write your data →	Last name:	First name:	ID number:	Points
EXERCISE 1				
EXERCISE 2				
TOTAL				

EXERCISE 1: Multiple-choice questions (20 points)

Remark: Only one choice is correct. Use the enclosed grid to select your choice. A correct answer awards 3 points, while a wrong answer awards a -1 penalization. The final result will be given by summing up all the obtained points (0 for a missing answer), by normalizing on a 20 base.

- Let $f(n)$ and $g(n)$ denote the message complexity of the *Chang & Roberts* algorithm in the average and in the worst case, respectively. Which of the following asymptotic relations is wrong?
*a) $f(n) = \Theta(g(n))$ b) $f(n) = O(g(n))$ c) $f(n) = o(g(n))$ d) $g(n) = \Omega(g(n))$
- Specify the largest among the following classes of rings for which the *leader election* problem can be solved through the *Hirschberger & Sinclair* algorithm:
a) asynchronous, anonymous, uniform, no-synchronized start b) synchronous, non-anonymous, uniform, no-synchronized start
c) asynchronous, non-anonymous, uniform, synchronized start *d) asynchronous, non-anonymous, uniform, no-synchronized start
- The most efficient *leader election* algorithm for a synchronous ring with n processors, non-anonymous and uniform, with minimum id m , has a message complexity of:
a) n b) it does not exist c) $\Theta(n \cdot m)$ *d) $\Theta(n)$
- Let us consider the *leader election* algorithm for a synchronous ring with n processors, non-anonymous and uniform. Let the minimum id in the ring be equal to 2^n . Then, the algorithm has a number of rounds of:
a) $O(n \cdot 2^n)$ b) $O(1)$ *c) $O(n \cdot 2^{2^n})$ d) $\Theta(n)$
- Let $f(n)$ and $g(n)$ denote the message complexity of the asynchronous versions of the *Prim* and the *GHS* algorithm, respectively, when executed on a dense graph, i.e., with $m = \Theta(n^2)$. Which of the following asymptotic relations is correct?
a) $f(n) = \Theta(g(n) \cdot n)$ *b) $f(n) = \Theta(g(n))$ c) $f(n) = o(g(n))$ d) $f(n) = \omega(g(n))$
- Let $f(n)$ and $g(n)$ denote the asymptotic number of rounds of the synchronous versions of the *Prim* and the *GHS* algorithm, respectively. Which of the following asymptotic relations is wrong?
a) $g(n) = O(f(n))$ *b) $f(n) = \Theta(g(n))$ c) $f(n) = \Omega(g(n))$ d) $f(n) = \omega(g(n))$
- The randomized algorithm for finding a *maximal independent set* running on a graph with n nodes and with degree $\Theta(\sqrt{n})$, with high probability has a number of phases in the order of:
*a) $O(\sqrt{n} \log n)$ b) $O(1)$ c) $O(\sqrt{n})$ d) $\Theta(n \log n)$
- Let G be an n -vertex graph of degree Δ . What is the approximation ratio guaranteed by the greedy algorithm for the *minimum dominating set* problem?
*a) $H(\Delta + 1)$ b) $H(\ln \Delta + 1)$ c) $\ln(H(\Delta))$ d) Δ
- In the *bakery algorithm* with n processors, which of the following is the second semaphore of the entry section of p_i ?
a) wait until $\text{Choosing}[j] = \text{false}$ or $(\text{Number}[j], j) > (\text{Number}[i], i)$ *b) wait until $\text{Number}[j] = 0$ or $(\text{Number}[j], j) > (\text{Number}[i], i)$
c) wait until $\text{Choosing}[j] = \text{false}$ d) wait until $\text{Number}[j] = 0$ or $(\text{Number}[j], j) < (\text{Number}[i], i)$
- In the bounded-space 2-processor *Mutex algorithm* with no lockout, which of the following is the first semaphore of the entry section of p_i ?
*a) wait until $\text{W}[1 - i] = 0$ or $\text{Priority} = i$ b) wait until $\text{W}[1 - i] = 1$ or $\text{Priority} = i$
c) wait until $\text{W}[1 - i] = 0$ or $\text{Priority} = 1 - i$ d) wait until $\text{Priority} = i$

Answer Grid

	Question									
Choice	1	2	3	4	5	6	7	8	9	10
a										
b										
c										
d										

EXERCISE 2: Open questions (10 points)

Remark: Select at your choice one out of the following two questions, and address it exhaustively.

- Describe and analyze the *slow-fast message* algorithm for the leader election problem.
- Describe and analyze the synchronous version of the *GHS* algorithm for the minimum spanning tree problem.